Content-Based Distributed Event-Based System

Master CILS — Module IAAIO, Revision : 1647

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October 2017
Most of the content of these slides is extracted from the following references:

Outline

1. Content-based data and filter models
2. Distributed Event-Based Systems
3. Formal specification of distributed routing
4. Distributed routing algorithms
5. Lab on content-based filtering with the muDEBS framework
1 Content-based data and filter models

1.1 Data model and Filter model
1.2 Structured records
1.3 Semi-structured records
1.4 Objects
1.1 Data model and Filter model

- **Data model:** how the content of notifications is structured
- **Filter model:** how subscriptions can be specified
  - How notifications can be selected by applying filters that evaluate predicates over the content of notifications
1.1.1 Tuples

Data model:

- A notification is a tuple: an ordered set of attributes

Filter model:

- A subscription is defined as a template
- The attributes of notifications and templates are matched to each other according to their position

Example: the notification \((\text{StockQuote}, \text{“Foo Inc”}, 45)\) is matched by the subscription template \((\text{StockQuote}, \text{“Foo Inc”}, *)\)

- Tuples with templates provide a simple model that is not flexible
  - Because attributes cannot be optional
1.2 Structured records
1.2.1 Data model

- A notification $n$ is a nonempty set of attributes $\{a_1, ..., a_n\}$
- $a_i$ is a (name,value) pair: $(n_i, v_i)$
- Attribute names are unique: $i \neq j \Rightarrow n_i \neq n_j$
- Example of notification: $\{(type, StockQuote),(name, “Infineon”), (price, 45.0)\}$
- More powerful than tuples since attributes can be optional in subscriptions and notifications
1.2.2 Filter model

- **Attribute filter**: triple $A_i = (n_i, Op_i, C_i)$
  with $n_i =$ attribute name, $Op_i =$ test operator, $C_i =$ value for the test

- **$L_A(A_i)$**: set of values $v_i$ that cause an attribute filter to match attribute $n_i$
  - $L_A(A_i) = \{v_i| Op_i(v_i, C_i) = true\}$
  - Usually $L_A(A_i) \neq \emptyset$

- **Filter $F$**: boolean function applied to a notification $n$: $F(n) \rightarrow \{true, false\}$

- **The set of matching notifications $N(F) = \{n| F(n) = true\} \subseteq \mathcal{N}$$

- **Simple filter** = filter consisting of a single atomic predicate

- **Compound filter** = conjunction of simple filters: $F = A_1 \wedge ... \wedge A_n$
  - E.g., $(type = \text{StockQuote}) \wedge (name = \text{“Foo Inc”}) \wedge (price \notin [30, 40])$

- **Notification $n$ matches a filter $F$ iff** it satisfies all the attributes filters of $F$
  - Attributes can be optional in the notification
  - New attributes can be added without affecting existing filters
1.2.3 Identity, overlapping, covering of attribute filters

- **Identity:**
  - \( A_1 \equiv A_2 \iff n_1 = n_2 \land L_A(A_1) = L_A(A_2) \)
  - *E.g.*, \((price \in \{20, 21, 22, 23, 24, 25\})\) is identical to \((price \in [20, 25])\)

- **Overlapping:**
  - \( A_1 \cap A_2 \iff n_1 = n_2 \land L_A(A_1) \cap L_A(A_2) \neq \emptyset \)
  - *E.g.*, \((price > 25)\) overlaps \((price \in [20, 30])\)

- **Covering:**
  - \( A_1 \sqsupseteq A_2 \iff n_1 = n_2 \land L_A(A_1) \supseteq L_A(A_2) \)
  - *E.g.*, \(A_1 = (price > 10)\) covers \(A_2 = (price \in [20, 30])\)

- **Disjoint**
  - \( A_1 \nsubseteq A_2 \iff n_1 = n_2 \land L_A(A_1) \cap L_A(A_2) = \emptyset \)
  - \{\(price < 10\)\} and \{\(price > 20\)\} are disjoint
1.2.4 Routing optimisations with identity

An identity test among filters is necessary to implement identity-based routing to avoid redundant routing entries and unnecessary forwarding of (un)subscriptions.

Given two filters $F_1 = A_1^1 \land ... \land A_n^1$ and $F_2 = A_1^2 \land ... \land A_m^2$ that are conjunctions of attribute filters with at most one attribute filter per attribute, $F_1 \equiv F_2$ iff they contain the same number of attributes filters $\land (\forall i, \exists j : A_i^1 \equiv A_j^2)$

E.g., $F_1 = \{x = 4\} \land \{y > 5\}$ not identical to $F_2 = \{x = 4\} \land \{y > 5\} \land \{z \in [3, 5]\}$
1.2.5 Routing optimisations with covering

- A covering test among filters is necessary to implement covering-based routing to avoid redundant routing entries and unnecessary forwarding of (un)subscriptions ∧ to get rid of the obsolete routing entries.

- Given two filters $F_1 = A_1^1 \land \ldots \land A_n^1$ and $F_2 = A_1^2 \land \ldots \land A_m^2$ that are conjunctions of attribute filters with at most one attribute filter per attribute,
  $$F_1 \sqsubseteq F_2 \iff \forall i, \exists j : A_i^1 \sqsubseteq A_j^2$$

- E.g., $F_1 = \{x = 4\} \land \{y > 5\}$ covers $F_2 = \{x = 4\} \land \{y > 5\} \land \{z \in [3, 5]\}$

- E.g., $F_3 = \{x \geq 2\} \land \{y > 5\}$ covers $F_4 = \{x = 4\} \land \{y = 7\} \land \{z \in [3, 5]\}$

---

1. A routing entry covers another routing entry, which becomes obsolete
1.2.6 Covering with types and comparison

- $n_1 = n_2$

- Covering among notification types
  - A notification $n$ is an instance of Type $T$: $n \text{ instanceof } T$

<table>
<thead>
<tr>
<th>$A_1$</th>
<th>$A_2$</th>
<th>$A_1 \sqsupseteq A_2 \text{ iff}$</th>
</tr>
</thead>
<tbody>
<tr>
<td>$n \text{ instanceof } T_1$</td>
<td>$n \text{ instanceof } T_2$</td>
<td>$T_1 = T_2 \lor T_1 \text{ supertypeof } T_2$</td>
</tr>
</tbody>
</table>

- Covering among comparison constraints on simple values

<table>
<thead>
<tr>
<th>$A_1$</th>
<th>$A_2$</th>
<th>$A_1 \sqsupseteq A_2 \text{ iff}$</th>
</tr>
</thead>
<tbody>
<tr>
<td>$x \neq c_1$</td>
<td>$x &lt; c_2$</td>
<td>$c_1 \geq c_2$</td>
</tr>
<tr>
<td>$x &gt; c_1$</td>
<td>$x &gt; c_2$</td>
<td>$c_1 \leq c_2$</td>
</tr>
</tbody>
</table>

- E.g., $A_1 = (x \neq 15)$ and $A_2 = (x < 10)$ $\implies A_1 \sqsupseteq A_2$
- E.g., $A_1 = (x > 10)$ and $A_2 = (x > 20)$ $\implies A_1 \sqsupseteq A_2$
1.2.7 Covering with intervals and strings

- \( n_1 = n_2 \)

- Covering among interval constraints on simple values

\[
\begin{array}{c|c|c}
A_1 & A_2 & A_1 \sqsupseteq A_2 \text{ iff} \\
\hline
x \in I_1 & x \in I_2 & l_1 \supseteq l_2 \\
\hline
x \notin I_1 & x \notin I_2 & l_1 \subseteq l_2 \\
\end{array}
\]

- E.g. \( A_1 = (x \in [3, 10]) \) and \( A_2 = (x \in [4, 6]) \) \( \implies A_1 \sqsupseteq A_2 \)

- Covering among constraints on strings

\[
\begin{array}{c|c|c}
A_1 & A_2 & A_1 \sqsupseteq A_2 \text{ iff} \\
\hline
s \text{ hasPrefix } S_1 & s \text{ hasPrefix } S_2 & S_2 \text{ hasPrefix } S_1 \\
\hline
s \text{ hasPostfix } S_1 & s \text{ hasPostfix } S_2 & S_2 \text{ hasPostfix } S_1 \\
\hline
s \text{ hasSubstring } S_1 & s \text{ hasSubstring } S_2 & S_2 \text{ hasSubstring } S_1 \\
\end{array}
\]

- E.g. \( A_1 = (s \text{ hasPrefix } "abc") \) and \( A_2 = (s \text{ hasPrefix } "abcd") \) \( \implies A_1 \sqsupseteq A_2 \)
1.2.8 Covering with sets

- $n_1 = n_2$

Covering among set constraints on simple values

<table>
<thead>
<tr>
<th>$A_1$</th>
<th>$A_2$</th>
<th>$A_1 \supseteq A_2$ iff</th>
</tr>
</thead>
<tbody>
<tr>
<td>$x \in M_1$</td>
<td>$x \in M_2$</td>
<td>$M_1 \supseteq M_2$</td>
</tr>
<tr>
<td>$x \notin M_1$</td>
<td>$x \notin M_2$</td>
<td>$M_1 \subseteq M_2$</td>
</tr>
</tbody>
</table>

Covering among set constraints on multi values

<table>
<thead>
<tr>
<th>$A_1$</th>
<th>$A_2$</th>
<th>$A_1 \supseteq A_2$ iff</th>
</tr>
</thead>
<tbody>
<tr>
<td>$X \subset M_1$</td>
<td>$X \subset M_2$</td>
<td>$M_1 \supset M_2$</td>
</tr>
<tr>
<td>$X \text{ contains } a_1$</td>
<td>$X \supset M_2$</td>
<td>$a_1 \in M_2$</td>
</tr>
<tr>
<td>$X \supset M_1$</td>
<td>$X \supset M_2$</td>
<td>$M_1 \subset M_2$</td>
</tr>
<tr>
<td>$X \text{ notContains } a_1$</td>
<td>$X \text{ disjunct } M_2$</td>
<td>$a_1 \in M_2$</td>
</tr>
<tr>
<td>$X \text{ disjunct } M_1$</td>
<td>$X \text{ disjunct } M_2$</td>
<td>$M_1 \subset M_2$</td>
</tr>
<tr>
<td>$X \text{ overlaps } M_1$</td>
<td>$X \text{ overlaps } M_2$</td>
<td>$M_1 \supset M_2$</td>
</tr>
</tbody>
</table>
1.2.9 Routing optimisations with overlapping

- An overlapping test among filters is necessary to use advertisements in subscription-based routing optimisations.

- Advertisement and subscription routing tables are used to route (un)subscriptions from consumers to producers.
  - A subscription can be served by an advertisement if both overlap.

- Given two filters \( F_1 = A_1^1 \land ... \land A_n^1 \) and \( F_2 = A_1^2 \land ... \land A_m^2 \) that are conjunctions of attribute filters with at most one attribute filter per attribute,
  - \( F_1 \) and \( F_2 \) are disjoint iff \( \exists i, j : (n_i^1 = n_j^2) \land (L_A(A_i^1) \cap L_A(A_j^2) = \emptyset) \)
    - E.g., \( F_1 = \{x \geq 2\} \land \{y > 5\} \) and \( F_2 = \{x < 1\} \land \{y < 7\} \) are disjoint because \( \{x \geq 2\} \) and \( \{x < 1\} \) are disjoint.
  - \( F_1 \) and \( F_2 \) overlap iff \( \forall i, j : (n_i^1 = n_j^2) \land (L_A(A_i^1) \cap L_A(A_j^2) = \emptyset) \)
    - E.g., \( F_1 = \{x \geq 2\} \land \{y > 5\} \) and \( F_2 = \{x < 5\} \land \{y < 7\} \) because \( \{x \geq 2\} \) overlaps \( \{x < 5\} \) and \( \{y > 5\} \) overlaps \( \{y < 7\} \)
1.2.10 Routing optimisations with merging

- Merging of conjunctive filters

- A merging test among filters is necessary to implement merging-based routing to reduce the number of subscriptions and advertisements stored by brokers

- Examples:
  - \( F_1 = \{ x = 5 \} \land \{ y \in \{2, 3\} \} \) and \( F_2 = \{ x = 5 \} \land \{ y \in \{4, 5\} \} \) can be merged to \( F = \{ x = 5 \} \land \{ y \in \{2, 3, 4, 5\} \} \)

  - \( F_1 = \{ y = 3 \} \land \{ x = 5 \} \) and \( F_2 = \{ y = 3 \} \land \{ x \neq 5 \} \) can be merged to \( F = \{ y = 3 \} \)

- Example of **perfect** merging rules for attribute filters

<table>
<thead>
<tr>
<th>( A_1 )</th>
<th>( A_2 )</th>
<th>Condition</th>
<th>( A_1 \cup A_2 )</th>
</tr>
</thead>
<tbody>
<tr>
<td>( x \in M_1 )</td>
<td>( x \in M_2 )</td>
<td>-</td>
<td>( x \in M_1 \cup M_2 )</td>
</tr>
<tr>
<td>( x \notin M_1 )</td>
<td>( x \notin M_2 )</td>
<td>( M_1 \cap M_2 = \emptyset ) ( M_1 \cap M_2 \neq \emptyset )</td>
<td>( \exists x ) (i.e., no att. filter) ( x \notin M_1 \cap M_2 )</td>
</tr>
<tr>
<td>( X \text{ overlaps } M_1 )</td>
<td>( X \text{ overlaps } M_2 )</td>
<td>-</td>
<td>( X \text{ overlaps } M_1 \cup M_2 )</td>
</tr>
<tr>
<td>( X \text{ disjunct } M_1 )</td>
<td>( X \text{ disjunct } M_2 )</td>
<td>( M_1 \cap M_2 = \emptyset )</td>
<td>( \exists X ) (i.e., no att. filter)</td>
</tr>
</tbody>
</table>

17/49 10/2017
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Content-Based Distributed Event-Based System
1.3 Semi-structured records
1.3.1 Data model

- Notification = XML document = set of elements arranged in a tree
  - Element = set of attributes + subordinate child elements
    - Attributes = pairs \((name, value)\)
    - Sibling attributes can have same name \(\implies\) names address sets of attr.

```xml
<notification>
  <auction endtime="05/18/02 22:17:42"
         minprice="50">
    <seller name="Smith" id="1234"/>
    <item>
      <board ...
    </item>
    <item>
      <cpu manufacturer="AMD"
           type="Athlon"
           clock="800"/>
    </item>
  </auction>
</notification>
```
1.3.2 Filter model I

- A filter model uses a path expression (e.g., XPath)
  - Select a set of attributes
  - Impose constraints on the selected attributes

- A filter is a conjunction of path filters: \( F = \land_i P_i \)

- A path filter \( P = (S, C) \) consists of an element selector \( S \) and an element filter \( C \)

- An element selector selects a subset of the elements of a notification
  - An absolute path: e.g. `/notification/auction/item/cpu`
  - An abbreviated path: e.g. `//cpu`

- An element filter is a conjunction of a nonempty set of attribute filters:
  \( C = \land_i A_i \)
  - e.g. `[@manufacturer = “AMD” \land @clock \geq 700]`
1.3.2 Filter model II

- Example of path filter:
  `/notification/auction/item/cpu[@manufacturer = “AMD” ∧ @clock ≥ 700]`

- $L_A(A)$: set of all values that cause an attribute filter $A$ to match an attribute

- $A_1 = (n_1, Q_1)$ covers $A_2 = (n_2, Q_2)$,
  $A_1 ⊒ A_2$ iff $n_1 = n_2$ ∧ $L_A(A_1) ⊇ L_A(A_2)$

  - Example: $[@clock ≥ 600]$ covers $[@clock ≥ 700]$

- $L_E(C)$: set of all elements that match an element filter $C$

- $C_1$ covers $C_2$, $C_1 ⊒ C_2$ iff $L_E(C_1) ⊇ L_E(C_2)$

  - Example: $[@clock ≥ 600]$ covers $[@manufacture = “AMD” ∧ @clock ≥ 700]$

- $C_1$ is disjoint with $C_2$ if there exists no attribute that is constrained in both element filters

  - Example: $[@minprice ¡ 100]$ is disjoint with $[@name = “Pu”]$
1.3.2 Filter model III

- $L_S(S)$: set of all elements that are selected by an element selector $S$
- $S_1$ covers $S_2$, $S_1 \supseteq S_2$ iff $L_S(S_1) \supseteq L_S(S_2)$
- $S_1$ is disjoint with $S_2$ iff $L_S(S_1) \cap L_S(S_2) = \emptyset$
- An absolute path covers another absolute path iff both are identical
- An abbreviated path covers another (abbreviated/absolute) path iff the former is a suffix of the later
  
  **Example:** //cpu covers //item/cpu because //cpu selects all elements named cpu, //item/cpu only selects those elements named cpu which are a sub-element of an element item

- $L_P(P)$: set of all elements that match a path filter $P$
1.3.2 Filter model IV

- $P_1 = (S_1, C_1)$ covers $P_2 = (S_2, C_2)$, $P_1 \sqsupseteq P_2$ iff $L_P(P_1) \supseteq L_P(P_2)$

  - Example: `//cpu[@manufacturer = "AMD"]` covers `//cpu[@manufacturer = "AMD" ∧ @clock ≥ 700]`

- $P_1$ is disjoint with $P_2$ iff $S_1$ is disjoint with $S_2$ or $C_1$ is disjoint with $C_2$

- Lemma: Given two path filters $P_1 = (S_1, C_1)$ and $P_2 = (S_2, C_2)$:
  $P_1 \sqsupseteq P_2$ iff $S_1 \sqsupseteq S_2$ ∧ $C_1 \sqsupseteq C_2$
  A filter $F_1$ covers $F_2$, $F_1 \sqsupseteq F_2$ iff $N(F_1) \supseteq N(F_2)$

- Lemma: Given two filters $F_1 = P_1^1 ∧ ... ∧ P_n^1$ and $F_2 = P_1^2 ∧ ... ∧ P_m^2$:
  $F_1 \sqsupseteq F_2$ iff $\forall i \exists j$ such that $P_i^1 \sqsupseteq P_j^2$

  - Example: the filter `{//cpu[@type = "Athlon"]}` covers `{//seller[@name = "Pu"] ∧ //cpu[@type = "Athlon" ∧ @clock ≥ 600]}`
1.4 Objects

- Model notifications and filters as objects
- Calling methods on attribute objects
  - Methods can be invoked on the objects embedded in the notification
  - The return value of the method can be a boolean value that is interpreted as a result of the attribute filter or a value that is used to evaluate the constraint

  **Example:** An instance of a class *StockQuote* has been embedded in a notification
  - The object possesses an attribute with the name *quote*
  - \( A \equiv (quote.id() = \text{“IBM”}) \)
  - \( A \) covers
    \( (quote.isRealTime()) \land (quote.id() = \text{“IBM”}) \land (quote.price() > 45.0) \)
2 Distributed Event-Based Systems

2.1 Distributed system model
2.2 Distributed notification service architecture
2.3 Distributed notification routing
2.1 Distributed system model

- Each node runs one or more processes
- Processes interact by passing messages via links between them
- A link connects a pair of processes and transmits messages asynchronously
- A FIFO ordering of messages is applied
2.2 Distributed notification service architecture

- The notification service forms an overlay network in the underlying system
- The overlay consists of event brokers that run as processes on physical nodes
  - Border and inner brokers forward the message to neighbouring brokers according to filter-based routing tables and routing strategies
  - Messages are sent to local brokers
  - Local brokers deliver the message to the application components
2.3 Distributed notification routing

Routing =

- Matching of all the notifications with all the subscriptions
- Delivering of the notifications to all the clients and the neighbouring brokers with a matching subscription

1. Flooding:

- Brokers forward notifications to all the neighbouring brokers
- Only brokers to which subscribers are connected test on matching subscriptions
- Advantage: guarantee that all the notifications will reach their destination
- Drawback: many unnecessary messages are exchanged among brokers

2. Filter-based: depends on routing tables (RT), which are maintained by brokers

- A routing entry is a filter-destination pair \((F, D)\)
- Entries are updated by sending control messages

3. Gossiping (not studied here)
2.3.1 Strategy of filter-based routing

- **Simple routing:**
  - Each broker has global knowledge about all active subscriptions
  - Routing tables may grow excessively
  - High filter forwarding overhead if subscriptions change frequently

- **Advanced routing:**
  - **Identity-based routing:** avoids forwarding of a subscription that matches identical subscriptions
  - **Covering-based routing:** avoids forwarding of those subscriptions that only accept a subset of notifications matched by a previously forwarded subscription
  - **Merging-based routing:** the broker creates a new cover for the merged routing entries that replaces the old ones
3 Formal specification of distributed routing

3.1 Distributed system model for routing
3.2 Notations for notif. forwarding and delivery
3.3 Valid routing
3.4 Safety and liveness conditions of valid routing
3.5 Safety and liveness conditions of monotone valid routing
3.1 Distributed system model for routing

Assumptions:

- Acyclic connected topologies
- Topology is static and clients are stationary
- Communication channels are reliable and respect FIFO message ordering
- Message delay is unknown but finite, and system is not overloaded and fault-free
- System without advertisements
3.2 Notations for notif. forwarding and delivery

- \( T_B^D \): set of filters of the routing table of broker \( B \) regarding single destination \( D \)
  \[
  T_B^D = \{F|\exists(F, D) \in T_B\}
  \]

- \( T_B^\backslash D \): set of filters regarding all but single destination \( D \)
  \[
  T_B^\backslash D = \{F|\exists(F, E) \in T_B \land E \neq D\}
  \]

- \( N(T_B^D) \): set of notifications that match \( T_B^D \)

- \( N_B \): set of neighbouring brokers

- \( L_B \): set of local consumers

- \( F_B(n) \): destinations to which a broker \( B \) forwards or delivers a notification \( n \)
  \[
  F_B(n) = \{D|D \in N_B \cup L_B \land n \in N(T_B^D)\}
  \]
3.3 Valid routing

- Valid routing algorithm = adapts the routing configuration by preserving the safety and liveness properties of the DEBS

- Additional notations:
  - $\theta(Y)$: identity of the broker that manages consumer $Y$
  - Simple directed path connecting a broker with $\theta(Y)$—i.e., the access broker
    - $B_1, \ldots, B_j$ simple path in the network of brokers
    - $\gamma(B_1, \ldots, B_j)$: set of notifications such that if a notification is published at $B_j$ and stays in this set, it reaches $B_1$ over this path
      - $\gamma(B_1, \ldots, B_j) = \bigcap_{1 < k \leq j} N(T_{B_k}^{B_{k-1}})$
3.4 Safety and liveness conditions of valid routing

- To guarantee safety, the local routing configuration ensures that only matching notifications are delivered
  
  • Local subset validity

  \[ \square \left[ N(T^Y_{\theta(Y)}) \subseteq N(S_Y) \right] \]  
  
  \( (=r1) \)

- To guarantee liveness, when a consumer \( Y \) subscribes to a filter \( F \) and stays subscribed, then from some time, every notification that is published at any broker \( B \) and that matches \( F \) should be delivered to \( Y \)
  
  • Eventual super-set validity

  \[ \square \left[ \square (F \in S_Y) \implies \diamond \square \left[ N(T^Y_{\theta(Y)}) \supseteq N(F) \right] \right] \]  
  
  \( (=r2: \text{From } \theta(Y) \text{ to } Y) \)

  \[ \square \left[ \square (F \in S_Y) \land B \neq \theta(Y) \land n \in N(F) \right] \implies \diamond \square \left[ n \in \gamma(\theta(Y),...,B) \right] \]  
  
  \( (=r3: \text{From } B \text{ to } \theta(Y)) \)
3.4.1 Monotone valid routing algorithms

- Drawbacks of valid routing
  - Local subset validity does not require immediate delivery
  - Eventual super-set validity is a property of the routing configuration of the entire topology

- Improvements
  - Immediate delivery
    - Local consumer subscription followed by local publisher publication should imply local notification of the consumer
  - Set of notifications forwarded is monotonically increasing for any path
    - Notifications sent over $B_{i+1} \rightarrow B_i$ are sent over $B_{i+2} \rightarrow B_{i+1}$
      - Only depends on the routing configurations of neighbouring brokers
3.5 Safety and liveness conditions of monotone valid routing

- Reminder:
  - $T_B^D = \{ F | \exists (F, D) \in T_B \}$
  - $T_B^\ \ D = \{ F | \exists (F, E) \in T_B \land E \neq D \}$

- Local validity $\equiv$ immediate delivery

  $\square \left[ N(T_{\theta(Y)}^Y) = N(S_Y) \right] \quad (= \text{merging of r1 and r2 + strengthness})$

- Eventual monotone remote validity

  $\square \left[ \square [ n \in N(T_{B_i}^{B_j}) ] \implies \Diamond \square [ n \in N(T_{B_j}^{B_i}) ] \right]$

2. If $n$ is forwarded to $B_k \neq B_j \in N_{B_i}$ then $n$ comes from $B_j$. 
4 Distributed routing algorithms

4.1 Generic algorithm
4.2 Flooding
4.3 Simple routing
4.4 Identity-based routing
4.1 Generic algorithm
4.1.1 Main program

The main program starts when the broker is created:

1. Initialise the routing table $T_B$ of the broker $B$
2. Initialise a delivery queue $Q_C$ for each local consumer $C$
3. Enter an infinite loop that dispatches messages arriving from neighbouring brokers to the `handleMessage` procedure

```
Program ContentBasedRouting()
initialise $T_B$
initialise $Q_C$ for all $C \in L_B$
loop
wait until a message is available
$m \leftarrow$ next selected message
handleMessage($m$)
```
4.1.2 handleMessage procedure

handleMessage dispatches a message based on message type

- Two types of messages are exchanged among neighbouring brokers
  1. forward($n$): to disseminate a notification $n$ in the network of brokers
  2. admin($S$, $U$): to propagate routing table updates
     - $S$: set of subscriptions
     - $U$: set of unsubscriptions
  3. administer($S$, $U$): to compute the admin messages to send
     - $M_S$: set of pairs ($\text{filter}_{\text{sub}}$, destination) for sending admin messages
     - $M_U$: set of pairs ($\text{filter}_{\text{unsub}}$, destination) for sending admin messages

```plaintext
1 procedure handleMessage(Message m)
2   if m is forward($n$) from neighbour $u$ then
3     handleNotification($u$, $n$)
4   if m is admin($S$, $U$) from neigh. $u$ then
5     ($M_S$, $M_U$) ← administer($u$, $S$, $U$)
6     handleAdminMessage($u$, $M_S$, $M_U$)
```
4.1.3 handleNotification procedure

- handleNotification sends forward messages to neighbouring brokers
- handleNotification notifies local consumers
  - notify is called by the broker to notify a local consumer about a notification
  - The notification is appended to the delivery queue $Q_Y$ of the consumer $Y$

1. procedure
   handleNotification(Neighbour $D$, Notification $n$)
2. send "forward($n$)" to all the neighbours $\in F_B(n) \setminus \{D\}$
3. forall local consumers $C \in F_B(n)$ do
   notify($C$, $n$)
5. procedure notify(Consumer $Y$, Notification $n$)
6. $Q_Y \leftarrow \text{append}(Q_Y, n)$
4.1.4 Preliminary words about the generic administer procedure

- The code of `administer` is implemented by framework instantiations to realise a concrete routing algorithm
  - Flooding
  - Simple
  - Identity-based
  - Covering-based
  - Perfect merging
  - Imperfect merging

- `administer` returns two sets that are pairs: \((\text{filter}_{\text{sub}}, \text{destination})\) or \((\text{filter}_{\text{unsub}}, \text{destination})\)

- Send an admin message to destination for \(\text{filter}_{\text{sub}}\)
  - Sending done in `handleAdminMessage`, as explained in next slide
4.1.5 handleAdminMessage procedure

- The values returned by `administer` are used as input to `handleAdminMessage`

- `handleAdminMessage` sends admin messages to neighbouring brokers

```
procedure handleAdminMessage(Dest D, Set M_S, Set M_U)

forall B_i ∈ N_B \{D\}

S' ← {F| (F, B_i) ∈ M_S}
U' ← {F| (F, B_i) ∈ M_U}

if S' ≠ ∅ ∨ U' ≠ ∅ then
    send "admin(S', U')" to B_i
```

1. admin (S, U)
2. admin (S', U')
4.1.6 pub, sub and unsub procedures

- pub is called by a local publisher to publish a notification
- sub is called by a local consumer to subscribe to a filter
- unsub is called by a local consumer to unsubscribe to a filter

```plaintext
procedure pub (Publisher X, Notification n)
  handleNotification(X, n)

procedure sub (Consumer Y, Filter F)
  (MS, MU) ← administer(Y, {F}, ∅)
  handleAdminMessage(Y, MS, MU)

procedure unsub (Consumer Y, Filter F)
  (MS, MU) ← administer(Y, ∅, {F})
  handleAdminMessage(Y, MS, MU)
```

![Diagram](image)
4.2 Flooding

- Idea: a broker forwards a notification to all its neighbours
  - Each broker is initialised to the set \( \{(F_T, U) | U \in N_B \} \) with \( \forall n \in N, F_T(n) = true \)
- Each broker updates its routing table (RT) regarding its local consumers
  - If a consumer \( Y \) subscribes to a filter \( F \), the broker adds \( (F, Y) \) to its RT
  - If a consumer \( Y \) unsubscribes to a filter \( F \), the broker deletes \( (F, Y) \) from its RT
- Flooding does not require the remote routing configuration to be updated

```
1 procedure administer(Dest s, Set S, Set U)
2   \( T_B \leftarrow T_B \cup \{(F, s) | F \in S \} \)
3   \( T_B \leftarrow T_B \setminus \{(F, s) | F \in U \} \)
4   return (\emptyset, \emptyset);
```
4.3 Simple routing

- Idea: use filter forwarding to update the routing configuration in reaction to subscribing and unsubscribing consumers

- Initially, \( \forall B, T_B = \emptyset \)

1. **procedure** `administer`(Dest \( D \), Set \( S \), Set \( U \))
2. \( T_B \leftarrow T_B \cup \{(F, D) | F \in S\} \)
3. \( T_B \leftarrow T_B \setminus \{(F, D) | F \in U\} \)
4. \( M_S \leftarrow \{(F, H) | H \in N_B \setminus \{D\} \land F \in S\} \)
5. \( M_U \leftarrow \{(F, H) | H \in N_B \setminus \{D\} \land F \in U\} \)
6. return \((M_S, M_U)\)
4.4 Identity-based routing

Reminder:

- $T_B^D = \{ F | \exists (F, D) \in T_B \}$
- $T_B^\setminus D = \{ F | \exists (F, E) \in T_B \land E \neq D \}$

Idea: a subscription (unsubscription) is only forwarded to a neighbour $H$ if there is no identical subscription in the RT for a destination distinct from $H$

The superscript stands for *Identical*

- $C_B^{I}(F, D)$: set of routing entries in $T_B$ of which the filter is identical to the filter $F$ and of which the destination equals the destination $D$
  - $C_B^{I}(F, D) = \{ (G, D) | (G, D) \in T_B \land F \equiv G \}$

- $D_B^{I}(F)$: set of neighbours $H$ for which there is no routing entry $(G, D)$ in $T_B$, where $G$ is identical to $F$ and $D$ is distinct from $H$
  - $D_B^{I}(F) = \{ H \in N_B | \nexists G \in T_B^\setminus H : F \equiv G \}$
4.4.1 Algorithm

- If a broker $B$ receives a(n) (un)subscription from a neighbour or a consumer $D$: 
  - $B$ updates its RT (lines 4-6):
    - If $D$ is a neighbour, $B$ removes $C^l_B(F, D)$ (line 5)
    - If $D$ is a local consumer, $B$ removes solely $(F, D)$ (line 6)
  - $B$ forwards $F$ to all neighbours that are in $D^l_B(F)$ except $D$ (lines 7–10 and 13)
  - If $F$ is a subscription, $B$ inserts a routing entry $(F, D)$ into its RT (line 11)

```latex
\begin{algorithm}
\begin{algorithmic}
\Procedure{administer}{Dest $D$, Set $S$, Set $U$}
\State $M_S \leftarrow \emptyset$;
\State $M_U \leftarrow \emptyset$;
\ForAll{$F \in S \cup U$}
  \If{$D \in N_B$}
    \State $T_B \leftarrow T_B \setminus C^l_B(F, D)$;
  \Else
    \State $T_B \leftarrow T_B \setminus (F, D)$;
  \EndIf
\EndFor
\State $A \leftarrow \{ (F, H) | H \in D^l_B(F) \setminus \{D\} \}$;
\If{$F \in U$}
  \State $M_U \leftarrow M_U \cup A$;
\Else
  \State $M_S \leftarrow M_S \cup A$;
  \State $T_B \leftarrow T_B \cup \{ (F, D) \}$;
\EndIf
\State Return $(M_S, M_U)$;
\EndProcedure
\end{algorithmic}
\end{algorithm}
```
5 Lab on content-based filtering with the muDEBS framework

- A result of our research on context data distribution for the Internet of Things

  - **muDEBS**
    - muDEBS uses muSCA
    - muDEBS is used by muContext
    - muDEBS is used by QoCIM

- We propose to follow a tutorial on muDEBS in order to see in action the generic algorithm in the case “simple routing”